

```

(*****
(* HaskProofToStrong: convert HaskProof to HaskStrong *)
(*****)

```

Generalizable All Variables.

```

Require Import Preamble.
Require Import General.
Require Import NaturalDeduction.
Require Import Coq.Strings.String.
Require Import Coq.Lists.List.
Require Import Coq.Init.Specif.
Require Import HaskKinds.
Require Import HaskStrongTypes.
Require Import HaskStrong.
Require Import HaskProof.

```

Section HaskProofToStrong.

```

Context {VV:Type} {eqdec_vv:EqDecidable VV} {freshM:FreshMonad VV}.

```

```

Definition fresh := FMT_fresh freshM.
Definition FreshM := FMT freshM.
Definition FreshMon := FMT_Monad freshM.
Existing Instance FreshMon.

```

```

Definition ExprVarResolver  $\Gamma$  := VV -> LeveledHaskType  $\Gamma$  ★

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```

Definition judg2exprType (j:Judg) : Type :=
  match j with
  (  $\Gamma > \Delta > \Sigma \mid - \tau$  ) => forall ( $\xi$ :ExprVarResolver  $\Gamma$ ) vars,  $\Sigma = \text{mapOptionTree } \xi \text{ vars} \rightarrow$ 
    FreshM (ITree _ (fun t => Expr  $\Gamma \Delta \xi t$ )  $\tau$ )
  end.

```

```

Definition justOne  $\Gamma \Delta \xi \tau$  : ITree _ (fun t => Expr  $\Gamma \Delta \xi t$ ) [ $\tau$ ] -> Expr  $\Gamma \Delta \xi \tau$ .
  intros.
  inversion X; auto.
  Defined.

```

```

Definition ileaf '(it:ITree X F [t]) : F t.
  inversion it.
  apply X0.

```

Defined.

```
Lemma update_branches : forall  $\Gamma$  ( $\xi$ :VV  $\rightarrow$  LeveledHaskType  $\Gamma$   $\star$ ) lev l1 l2 q,  
  update_ $\xi$   $\xi$  lev (app l1 l2) q = update_ $\xi$  (update_ $\xi$   $\xi$  lev l2) lev l1 q.  
intros.  
induction l1.  
  reflexivity.  
  simpl.  
  destruct a; simpl.  
  rewrite IHl1.  
  reflexivity.  
Qed.
```

```
Lemma quark {T} (l1:list T) l2 vf :  
  (In vf (app l1 l2))  $\leftrightarrow$   
  (In vf l1)  $\wedge$  (In vf l2).
```

```
induction l1.  
simpl; auto.  
split; intro.  
right; auto.  
inversion H.  
inversion H0.  
auto.  
split.  
  
destruct IHl1.  
simpl in *.  
intro.  
destruct H1.  
left; left; auto.  
set (H H1) as q.  
destruct q.  
left; right; auto.  
right; auto.  
simpl.
```

```
destruct IHl1.  
simpl in *.  
intro.  
destruct H1.  
destruct H1.
```

```

left; auto.
right; apply H0; auto.
right; apply H0; auto.
Qed.

```

```

Lemma splitter {T} (l1:list T) l2 vf :
  (In vf (app l1 l2) →False)
  -> (In vf l1 →False) /\ (In vf l2 →False).
intros.
split; intros; apply H; rewrite quark.
auto.
auto.
Qed.

```

```

Lemma helper
: forall T Z {eqdt:EqDecidable T}(tl:Tree ??T)(vf:T) ξ (q:Z),
  (In vf (leaves tl) -> False) ->
  mapOptionTree (fun v' => if eqd_dec vf v' then q else ξ v') tl =
  mapOptionTree ξ tl.
intros.
induction tl;
  try destruct a;
  simpl in *.
set (eqd_dec vf t) as x in *.
destruct x.
subst.
  assert False.
  apply H.
  left; auto.
  inversion H0.
auto.
auto.
apply splitter in H.
destruct H.
rewrite (IHt11 H).
rewrite (IHt12 H0).
reflexivity.
Qed.

```

```

Lemma fresh_lemma'' Γ
: forall types ξ lev,

```

```

FreshM { varstypes : _
  | mapOptionTree (update_ξ(Γ:=Γ) ξ lev (leaves varstypes)) (mapOptionTree (@fst _ _) varstypes) = (types @@@ lev)
  /\ distinct (leaves (mapOptionTree (@fst _ _) varstypes)) }.
admit.
Defined.

Lemma fresh_lemma' Γ
: forall types vars Σ ξ lev, Σ = mapOptionTree ξ vars ->
FreshM { varstypes : _
  | mapOptionTree (update_ξ(Γ:=Γ) ξ lev (leaves varstypes)) vars = Σ
  /\ mapOptionTree (update_ξ ξ lev (leaves varstypes)) (mapOptionTree (@fst _ _) varstypes) = (types @@@ lev)
  /\ distinct (leaves (mapOptionTree (@fst _ _) varstypes)) }.
induction types.
intros; destruct a.
  refine (bind vf = fresh (leaves vars) ; return _).
  apply FreshMon.
  destruct vf as [ vf vf_pf ].
  exists [(vf,h)].
  split; auto.
  simpl.
  set (helper VV _ vars vf ξ (h@@lev) vf_pf) as q.
  rewrite q.
  symmetry; auto.
  simpl.
  destruct (eqd_dec vf vf); [ idtac | set (n (refl_equal _)) as n'; inversion n' ]; auto.
  split; auto.
  apply distinct_cons.
  intro.
  inversion H0.
  apply distinct_nil.
  refine (return _).
  exists []; auto.
  split.
  simpl.
  symmetry; auto.
  split.
  simpl.
  reflexivity.
  simpl.
  apply distinct_nil.
intros vars Σ ξ lev pf; refine (bind x2 = IHtypes2 vars Σ ξ lev pf; _).

```

```

apply FreshMon.
destruct x2 as [vt2 [pf21 [pf22 pfdist]]].
refine (bind x1 = IHtypes1 (vars,,(mapOptionTree (@fst _ _) vt2)) ( $\Sigma$ ,,(types2@@@1lev)) (update_ $\xi$   $\xi$  lev
  (leaves vt2)) _ _; return _).
apply FreshMon.
simpl.
rewrite pf21.
rewrite pf22.
reflexivity.
clear IHtypes1 IHtypes2.
destruct x1 as [vt1 [pf11 pf12]].
exists (vt1,,vt2); split; auto.

set (update_branches  $\Gamma$   $\xi$  lev (leaves vt1) (leaves vt2)) as q.
set (mapOptionTree_extensional _ _ q) as q'.
rewrite q'.
clear q' q.
inversion pf11.
reflexivity.

simpl.
set (update_branches  $\Gamma$   $\xi$  lev (leaves vt1) (leaves vt2)) as q.
set (mapOptionTree_extensional _ _ q) as q'.
rewrite q'.
rewrite q'.
clear q' q.
rewrite <- mapOptionTree_compose.
rewrite <- mapOptionTree_compose.
rewrite <- mapOptionTree_compose in *.
split.
destruct pf12.
rewrite H.
inversion pf11.
rewrite <- mapOptionTree_compose.
reflexivity.

admit.
Defined.

```

Lemma fresh_lemma Γ ξ vars Σ Σ' lev
: $\Sigma = \text{mapOptionTree } \xi \text{ vars } \rightarrow$

```

FreshM { vars' : _
  | mapOptionTree (update_ξ(Γ:=Γ) ξ lev ((vars',Σ')::nil)) vars = Σ
  /\ mapOptionTree (update_ξ ξ lev ((vars',Σ')::nil)) [vars'] = [Σ' @@ lev] }.
intros.
set (fresh_lemma' Γ [Σ'] vars Σ ξ lev H) as q.
refine (q >>>= fun q' => return _).
apply FreshMon.
clear q.
destruct q' as [varstypes [pf1 [pf2 pfdist]]].
destruct varstypes; try destruct o; try destruct p; simpl in *.
  destruct (eqd_dec v v); [ idtac | set (n (refl_equal _)) as n'; inversion n' ].
  inversion pf2; subst.
  exists v.
  destruct (eqd_dec v v); [ idtac | set (n (refl_equal _)) as n'; inversion n' ].
  split; auto.
  inversion pf2.
  inversion pf2.
Defined.

```

```

Definition ujudg2exprType Γ (ξ:ExprVarResolver Γ)(Δ:CoercionEnv Γ) Σ τ : Type :=
  forall vars, Σ = mapOptionTree ξ vars -> FreshM (ITree _ (fun t => Expr Γ Δ ξ t) τ).

```

```

Definition urule2expr : forall Γ Δ h j t (r:@Arrange _ h j) (ξ:VV -> LeveledHaskType Γ ★),
  ujudg2exprType Γ ξ Δ h t ->
  ujudg2exprType Γ ξ Δ j t
  .
intros Γ Δ.
  refine (fix urule2expr h j t (r:@Arrange _ h j) ξ {struct r} :
  ujudg2exprType Γ ξ Δ h t ->
  ujudg2exprType Γ ξ Δ j t :=
    match r as R in Arrange H C return
  ujudg2exprType Γ ξ Δ H t ->
  ujudg2exprType Γ ξ Δ C t

```

with

```

| RLeft   h c ctx r => let case_RLeft  := tt in (fun e => _) (urule2expr _ _ _ r)
| RRight  h c ctx r => let case_RRight := tt in (fun e => _) (urule2expr _ _ _ r)
| RCanL   a         => let case_RCanL  := tt in _
| RCanR   a         => let case_RCanR  := tt in _
| RuCanL  a         => let case_RuCanL := tt in _
| RuCanR  a         => let case_RuCanR := tt in _
| RAssoc  a b c     => let case_RAssoc := tt in _

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| RCossa a b c => let case_RCossa := tt in _
| RExch a b => let case_RExch := tt in _
| RWeak a => let case_RWeak := tt in _
| RCont a => let case_RCont := tt in _
| RComp a b c f g => let case_RComp := tt in (fun e1 e2 => _) (urule2expr _ _ _ f) (urule2expr _ _ _ g)
end); clear urule2expr; intros.

```

```

destruct case_RCanL.
simpl; unfold ujudg2exprType; intros.
simpl in X.
apply (X ([],,vars)).
simpl; rewrite <- H; auto.

```

```

destruct case_RCanR.
simpl; unfold ujudg2exprType; intros.
simpl in X.
apply (X (vars,,[])).
simpl; rewrite <- H; auto.

```

```

destruct case_RuCanL.
simpl; unfold ujudg2exprType; intros.
destruct vars; try destruct o; inversion H.
simpl in X.
apply (X vars2); auto.

```

```

destruct case_RuCanR.
simpl; unfold ujudg2exprType; intros.
destruct vars; try destruct o; inversion H.
simpl in X.
apply (X vars1); auto.

```

```

destruct case_RAassoc.
simpl; unfold ujudg2exprType; intros.
simpl in X.
destruct vars; try destruct o; inversion H.
destruct vars1; try destruct o; inversion H.
apply (X (vars1_1,,(vars1_2,,vars2))).
subst; auto.

```

```

destruct case_RCossa.
simpl; unfold ujudg2exprType; intros.

```

```
simpl in X.
destruct vars; try destruct o; inversion H.
destruct vars2; try destruct o; inversion H.
apply (X ((vars1,,vars2_1),,vars2_2)).
subst; auto.
```

```
destruct case_RExch.
simpl; unfold ujudg2exprType ; intros.
simpl in X.
destruct vars; try destruct o; inversion H.
apply (X (vars2,,vars1)).
inversion H; subst; auto.
```

```
destruct case_RWeak.
simpl; unfold ujudg2exprType; intros.
simpl in X.
apply (X []).
auto.
```

```
destruct case_RCont.
simpl; unfold ujudg2exprType ; intros.
simpl in X.
apply (X (vars,,vars)).
simpl.
rewrite <- H.
auto.
```

```
destruct case_RLeft.
intro vars; unfold ujudg2exprType; intro H.
destruct vars; try destruct o; inversion H.
apply (fun q => e  $\xi$  q vars2 H2).
clear r0 e H2.
simpl in X.
simpl.
unfold ujudg2exprType.
intros.
apply X with (vars:=vars1,,vars).
rewrite H0.
rewrite H1.
simpl.
reflexivity.
```



```

destruct case_RRight.
  intro vars; unfold ujudg2exprType; intro H.
  destruct vars; try destruct o; inversion H.
  apply (fun q => e ξ q vars1 H1).
  clear r0 e H2.
  simpl in X.
  simpl.
  unfold ujudg2exprType.
  intros.
  apply X with (vars:=vars,,vars2).
  rewrite H0.
  inversion H.
  simpl.
  reflexivity.

```

```

destruct case_RComp.
  apply e2.
  apply e1.
  apply X.
  Defined.

```

Definition letrec_helper $\Gamma \Delta l$ (varstypes:Tree ??(VV * HaskType $\Gamma \star$)) ξ' :

```

ITree (LeveledHaskType  $\Gamma \star$ )
  (fun t : LeveledHaskType  $\Gamma \star \Rightarrow$  Expr  $\Gamma \Delta \xi'$  t)
  (mapOptionTree ( $\xi' \circ$  (@fst _ _)) varstypes)
  -> ELetRecBindings  $\Gamma \Delta \xi'$  l varstypes.

```

```

intros.
induction varstypes.
destruct a; simpl in *.
destruct p.
simpl.
apply ileaf in X. simpl in X.
  apply ELR_leaf.
  rename h into  $\tau$ .
  destruct (eqd_dec (unlev ( $\xi'$  v))  $\tau$ ).
  rewrite <- e.
  destruct ( $\xi'$  v).
  simpl.
  destruct (eqd_dec h0 l).
  rewrite <- e0.

```

```

    apply X.
    apply (Prelude_error "level mismatch; should never happen").
    apply (Prelude_error "letrec type mismatch; should never happen").

apply ELR_nil.
apply ELR_branch.
  apply IHvarstypes1; inversion X; auto.
  apply IHvarstypes2; inversion X; auto.
Defined.

Definition unindex_tree {V}{F} : forall {t:Tree ??V}, ITree V F t -> Tree ??{ v:V & F v }.
  refine (fix rec t it := match it as IT return Tree ??{ v:V & F v } with
    | INone => T_Leaf None
    | ILeaf x y => T_Leaf (Some _)
    | IBranch _ _ b1 b2 => (rec _ b1),,(rec _ b2)
    end).
  exists x; auto.
Defined.

Definition fix_indexing X (F:X->Type)(J:X->Type)(t:Tree ??{ x:X & F x })
  : ITree { x:X & F x } (fun x => J (projT1 x))
  -> ITree X (fun x:X => J x) (mapOptionTree (@projT1 _ _) t).
  intro it.
  induction it; simpl in *.
  apply INone.
  apply ILeaf.
  apply f.
  simpl; apply IBranch; auto.
Defined.

Definition fix2 {X}{F} : Tree ??{ x:X & FreshM (F x) } -> Tree ??(FreshM { x:X & F x }).
  refine (fix rec t := match t with
    | T_Leaf None => T_Leaf None
    | T_Leaf (Some x) => T_Leaf (Some _)
    | T_Branch b1 b2 => T_Branch (rec b1) (rec b2)
    end).
  destruct x as [x fx].
  refine (bind fx' = fx ; return _).
  apply FreshMon.
  exists x.
  apply fx'.

```

Defined.

```
Definition case_helper tc  $\Gamma$   $\Delta$  lev tbranches avars  $\xi$  :
forall pcb:{sac : StrongAltCon & ProofCaseBranch tc  $\Gamma$   $\Delta$  lev tbranches avars sac},
  prod (judg2exprType (pcb_judg (projT2 pcb))) {vars' : Tree ??VV & pcb_freevars (projT2 pcb) = mapOptionTree  $\xi$  vars'} ->
  ((fun sac => FreshM
    { scb : StrongCaseBranchWithVVs VV eqdec_vv tc avars sac
      & Expr (sac_ $\Gamma$  sac  $\Gamma$ ) (sac_ $\Delta$  sac  $\Gamma$  avars (weakCK''  $\Delta$ )) (scbwv_ $\xi$  scb  $\xi$  lev) (weakLT' (tbranches @@ lev)) }) (projT1 pcb)).
  intro pcb.
  intro X.
  simpl in X.
  simpl.
  destruct pcb as [sac pcb].
  simpl in *.

  destruct X.
  destruct s as [vars vars_pf].

  refine (bind localvars = fresh_lemma' _ (unleaves (vec2list (sac_types sac _ avars)))) vars
    (mapOptionTree weakLT' (pcb_freevars pcb)) (weakLT'  $\circ$   $\xi$ ) (weakL' lev) _ ; _).
  apply FreshMon.
  rewrite vars_pf.
  rewrite <- mapOptionTree_compose.
  reflexivity.
  destruct localvars as [localvars [localvars_pf1 [localvars_pf2 localvars_dist ]]].
  set (mapOptionTree (@fst _ _) localvars) as localvars'.

  set (list2vec (leaves localvars')) as localvars''.
  cut (length (leaves localvars') = sac_numExprVars sac). intro H''.
  rewrite H'' in localvars''.
  cut (distinct (vec2list localvars'')). intro H'''.
  set (@Build_StrongCaseBranchWithVVs _ _ _ _ avars sac localvars'' H''') as scb.

  refine (bind q = (f (scbwv_ $\xi$  scb  $\xi$  lev) (vars,,(unleaves (vec2list (scbwv_exprvars scb)))))) _ ; return _).
  apply FreshMon.
  simpl.
  unfold scbwv_ $\xi$ .
  rewrite vars_pf.
  rewrite <- mapOptionTree_compose.
  clear localvars_pf1.
  simpl.
```

```

rewrite mapleaves'.

admit.

exists scb.
apply ileaf in q.
apply q.

admit.
admit.
Defined.

Definition gather_branch_variables
  Γ Δ (ξ:VV -> LeveledHaskType Γ ★ tc avars tbranches lev (alts:Tree ?? {sac : StrongAltCon &
    ProofCaseBranch tc Γ Δ lev tbranches avars sac})
  :
forall vars,
mapOptionTreeAndFlatten (fun x => pcb_freevars(Γ:=Γ) (projT2 x)) alts = mapOptionTree ξ vars
-> ITree Judg judg2exprType (mapOptionTree (fun x => pcb_judg (projT2 x)) alts)
-> ITree _ (fun q => prod (judg2exprType (pcb_judg (projT2 q)))
  { vars' : _ & pcb_freevars (projT2 q) = mapOptionTree ξ vars' })
alts.
  induction alts;
  intro vars;
  intro pf;
  intro source.
  destruct a; [ idtac | apply INone ].
  simpl in *.
  apply ileaf in source.
  apply ILeaf.
  destruct s as [sac pcb].
  simpl in *.
  split.
  intros.
  eapply source.
  apply H.
  clear source.

exists vars.
auto.

```

```

simpl in pf.
destruct vars; try destruct o; simpl in pf; inversion pf.
simpl in source.
inversion source.
subst.
apply IBranch.
apply (IHalts1 vars1 H0 X); auto.
apply (IHalts2 vars2 H1 X0); auto.

```

Defined.

Definition rule2expr : forall h j (r:Rule h j), ITree _ judg2exprType h -> ITree _ judg2exprType j.

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intros h j r.

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refine (match r as R in Rule H C return ITree _ judg2exprType H -> ITree _ judg2exprType C with
| RArrange a b c d e r      => let case_RURule := tt      in _
| RNote   Γ Δ Σ τ l n      => let case_RNote := tt      in _
| RLit    Γ Δ l            => let case_RLit := tt      in _
| RVar    Γ Δ σ            => let case_RVar := tt      in _
| RGlobal Γ Δ σ l wev      => let case_RGlobal := tt    in _
| RLam    Γ Δ Σ tx te x    => let case_RLam := tt      in _
| RCast   Γ Δ Σ σ τ γ x    => let case_RCast := tt      in _
| RAbsT   Γ Δ Σ κ σ a      => let case_RAbsT := tt      in _
| RAppT   Γ Δ Σ κ σ τ y    => let case_RAppT := tt      in _
| RAppCo  Γ Δ Σ κ σ1 σ2 γ σ l => let case_RAppCo := tt    in _
| RAbsCo  Γ Δ Σ κ σ σ1 σ2 y  => let case_RAbsCo := tt    in _
| RApp    Γ Δ Σ1 Σ2 tx te p  => let case_RApp := tt      in _
| RLet    Γ Δ Σ1 Σ2 σ1 σ2 p => let case_RLet := tt      in _
| RJoin   Γ p lri m x q    => let case_RJoin := tt in _
| RVoid   _ _              => let case_RVoid := tt in _
| RBrak   Σ a b c n m      => let case_RBrak := tt      in _
| REsc    Σ a b c n m      => let case_REsc := tt      in _
| RCase   Γ Δ lev tc Σ avars tbranches alts => let case_RCase := tt    in _
| RLetRec Γ Δ lri x y t    => let case_RLetRec := tt    in _
end); intro X_; try apply ileaf in X_; simpl in X_.

```

```

destruct case_RURule.

```

```

apply ILeaf. simpl. intros.
set (@urule2expr a b _ _ e r0 ξ) as q.

```

```

    set (fun z => q z) as q'.
    simpl in q'.
    apply q' with (vars:=vars).
    clear q' q.
    unfold ujudg2exprType.
    intros.
    apply X_ with (vars:=vars0).
    auto.
    auto.

destruct case_RBrak.
  apply ILeaf; simpl; intros; refine (X_ ξ vars H >>=> fun X => return ILeaf _ _). apply FreshMon.
  apply EBrak.
  apply (ileaf X).

destruct case_REsc.
  apply ILeaf; simpl; intros; refine (X_ ξ vars H >>=> fun X => return ILeaf _ _). apply FreshMon.
  apply EEsc.
  apply (ileaf X).

destruct case_RNote.
  apply ILeaf; simpl; intros; refine (X_ ξ vars H >>=> fun X => return ILeaf _ _). apply FreshMon.
  apply ENote; auto.
  apply (ileaf X).

destruct case_RLit.
  apply ILeaf; simpl; intros; refine (return ILeaf _ _).
  apply ELit.

destruct case_RVar.
  apply ILeaf; simpl; intros; refine (return ILeaf _ _).
  destruct vars; simpl in H; inversion H; destruct o. inversion H1. rewrite H2.
  apply EVar.
  inversion H.

destruct case_RGlobal.
  apply ILeaf; simpl; intros; refine (return ILeaf _ _).
  apply EGlobal.
  apply wev.

destruct case_RLam.

```

```

apply ILeaf.
simpl in *; intros.
refine (fresh_lemma _ ξ vars _ tx x H >>>= (fun pf => _)).
apply FreshMon.
destruct pf as [ vnew [ pf1 pf2 ]].
set (update_ξ ξ x ((vnew, tx )) :: nil)) as ξ' in *.
refine (X_ ξ' (vars,,[vnew]) _ >>>= _).
apply FreshMon.
simpl.
rewrite pf1.
rewrite <- pf2.
simpl.
reflexivity.
intro hyp.
refine (return _).
apply ILeaf.
apply ELam with (ev:=vnew).
apply ileaf in hyp.
simpl in hyp.
unfold ξ' in hyp.
apply hyp.

destruct case_RCast.
  apply ILeaf; simpl; intros; refine (X_ ξ vars H >>>= fun X => return ILeaf _ _). apply FreshMon.
  eapply ECast.
  apply x.
  apply ileaf in X. simpl in X.
  apply X.

destruct case_RJoin.
  apply ILeaf; simpl; intros.
  inversion X_.
  apply ileaf in X.
  apply ileaf in X0.
  simpl in *.
  destruct vars; inversion H.
  destruct o; inversion H3.
  refine (X ξ vars1 H3 >>>= fun X' => X0 ξ vars2 H4 >>>= fun X0' => return _).
  apply FreshMon.
  apply FreshMon.
  apply IBranch; auto.

```

```

destruct case_RApp.
  apply ILeaf.
  inversion X_.
  inversion X.
  inversion X0.
  simpl in *.
  intros.
  destruct vars. try destruct o; inversion H.
  simpl in H.
  inversion H.
  set (X1  $\xi$  vars1 H5) as q1.
  set (X2  $\xi$  vars2 H6) as q2.
  refine (q1 >>>= fun q1' => q2 >>>= fun q2' => return _).
  apply FreshMon.
  apply FreshMon.
  apply ILeaf.
  apply ileaf in q1'.
  apply ileaf in q2'.
  simpl in *.
  apply (EApp _ _ _ _ _ q1' q2').

destruct case_RLet.
  apply ILeaf.
  simpl in *; intros.
  destruct vars; try destruct o; inversion H.
  refine (fresh_lemma _  $\xi$  vars1 _  $\sigma_2$  p H1 >>>= (fun pf => _)).
  apply FreshMon.
  destruct pf as [ vnew [ pf1 pf2 ]].
  set (update_ $\xi$   $\xi$  p (((vnew,  $\sigma_2$  )) :: nil)) as  $\xi'$  in *.
  inversion X_.
  apply ileaf in X.
  apply ileaf in X0.
  simpl in *.
  refine (X  $\xi$  vars2 _ >>>= fun X0' => _).
  apply FreshMon.
  auto.

refine (X0  $\xi'$  (vars1,,[vnew]) _ >>>= fun X1' => _).
  apply FreshMon.
  rewrite H1.

```



```

simpl.
rewrite pf2.
rewrite pf1.
rewrite H1.
reflexivity.

refine (return _).
apply ILeaf.
apply ileaf in X0'.
apply ileaf in X1'.
simpl in *.
apply ELet with (ev:=vnew)(tv:=σ2).
apply X0'.
apply X1'.

destruct case_RVoid.
  apply ILeaf; simpl; intros.
  refine (return _).
  apply INone.

destruct case_RAppT.
  apply ILeaf; simpl; intros; refine (X_ ξ vars H >>>= fun X => return ILeaf _ _). apply FreshMon.
  apply ETyApp.
  apply (ileaf X).

destruct case_RAbsT.
  apply ILeaf; simpl; intros; refine (X_ (weakLT ∘ ξ) vars _ >>>= fun X => return ILeaf _ _). apply FreshMon.
  rewrite mapOptionTree_compose.
  rewrite <- H.
  reflexivity.
  apply ileaf in X. simpl in *.
  apply ETyLam.
  apply X.

destruct case_RAppCo.
  apply ILeaf; simpl; intros; refine (X_ ξ vars _ >>>= fun X => return ILeaf _ _). apply FreshMon.
  auto.
  eapply ECoApp.
  apply γ.
  apply (ileaf X).

```

```

destruct case_RAbsCo.
  apply ILeaf; simpl; intros; refine (X_ ξ vars _ >>>= fun X => return ILeaf _ _). apply FreshMon.
  auto.
  eapply ECoLam.
  apply (ileaf X).

destruct case_RLetRec.
  apply ILeaf; simpl; intros.
  refine (bind ξvars = fresh_lemma' _ y _ _ _ t H; _). apply FreshMon.
  destruct ξvars as [ varstypes [ pf1[ pf2 pfdist]]].
  refine (X_ ((update_ξ ξ t (leaves varstypes)))
    (vars,,(mapOptionTree (@fst _ _) varstypes)) _ >>>= fun X => return _); clear X_. apply FreshMon.
  simpl.
  rewrite pf2.
  rewrite pf1.
  auto.
  apply ILeaf.
  inversion X; subst; clear X.

  apply (@ELetRec _ _ _ _ _ _ varstypes).
  apply (@letrec_helper Γ Δ t varstypes).
  rewrite <- pf2 in X1.
  rewrite mapOptionTree_compose.
  apply X1.
  apply ileaf in X0.
  apply X0.

destruct case_RCCase.
  apply ILeaf; simpl; intros.
  inversion X_.
  clear X_.
  subst.
  apply ileaf in X0.
  simpl in X0.

(* body_freevars and alts_freevars are the types of variables in the body and alternatives (respectively) which are free
 * from the viewpoint just outside the case block -- i.e. not bound by any of the branches *)
rename Σ into body_freevars_types.
rename vars into all_freevars.
rename X0 into body_expr.
rename X into alts_exprs.

```

```

destruct all_freevars; try destruct o; inversion H.
rename all_freevars2 into body_freevars.
rename all_freevars1 into alts_freevars.

set (gather_branch_variables _ _ _ _ _ H1 alts_exprs) as q.
set (itmap (fun pcb alt_expr => case_helper tc  $\Gamma$   $\Delta$  lev tbranches avars  $\xi$  pcb alt_expr) q) as alts_exprs'.
apply fix_indexing in alts_exprs'.
simpl in alts_exprs'.
apply unindex_tree in alts_exprs'.
simpl in alts_exprs'.
apply fix2 in alts_exprs'.
apply treeM in alts_exprs'.

refine ( alts_exprs' >>>= fun Y =>
  body_expr  $\xi$  _ _
  >>>= fun X => return ILeaf _ (@ECase _ _ _ _ _ (ileaf X) Y)); auto.
  apply FreshMon.
  apply FreshMon.
  apply H2.
Defined.

```

```

Definition closed2expr : forall c (pn:@ClosedSIND _ Rule c), ITree _ judg2exprType c.
refine ((
  fix closed2expr' j (pn:@ClosedSIND _ Rule j) {struct pn} : ITree _ judg2exprType j :=
  match pn in @ClosedSIND _ _ J return ITree _ judg2exprType J with
  | cnd_weak          => let case_nil      := tt in INone _ _
  | cnd_rule h c cnd' r => let case_rule   := tt in rule2expr _ _ r (closed2expr' _ cnd')
  | cnd_branch _ _ c1 c2 => let case_branch := tt in IBranch _ _ (closed2expr' _ c1) (closed2expr' _ c2)
  end)); clear closed2expr'; intros; subst.
  Defined.

```

```

Lemma manyFresh : forall  $\Gamma$   $\Sigma$  ( $\xi$ 0:VV -> LeveledHaskType  $\Gamma$   $\star$ ),
  FreshM { vars : _ & {  $\xi$  : VV -> LeveledHaskType  $\Gamma$   $\star$  &  $\Sigma$  = mapOptionTree  $\xi$  vars } }.
intros  $\Gamma$   $\Sigma$ .
induction  $\Sigma$ ; intro  $\xi$ .
destruct a.
destruct l as [ $\tau$  l].
set (fresh_lemma'  $\Gamma$  [ $\tau$ ] [] []  $\xi$  l (refl_equal _)) as q.
refine (q >>>= fun q' => return _).
apply FreshMon.

```

```

clear q.
destruct q' as [varstypes [pf1 [pf2 distpf]]].
exists (mapOptionTree (@fst _ _) varstypes).
exists (update_ξ ξ 1 (leaves varstypes)).
symmetry; auto.
refine (return _).
exists [].
exists ξ; auto.
refine (bind f1 = IHΣ1 ξ ; _).
apply FreshMon.
destruct f1 as [vars1 [ξ1 pf1]].
refine (bind f2 = IHΣ2 ξ1 ; _).
apply FreshMon.
destruct f2 as [vars2 [ξ2 pf22]].
refine (return _).
exists (vars1,,vars2).
exists ξ2.
simpl.
rewrite pf22.
rewrite pf1.
admit.
Defined.

```

Definition proof2expr $\Gamma \Delta \tau \Sigma$ ($\xi_0: VV \rightarrow \text{LeveledHaskType } \Gamma$ ★

```

{zz:ToString VV} : ND Rule [] [ $\Gamma > \Delta > \Sigma \mid - [\tau]$ ] ->
FreshM (???{ ξ : _ & Expr  $\Gamma \Delta \xi \tau$ }).
intro pf.
set (closedFromSIND _ _ (mkSIND systemfc_all_rules_one_conclusion _ _ _ pf (scnd_weak [])) cnd_weak) as cnd.
apply closed2expr in cnd.
apply ileaf in cnd.
simpl in *.
clear pf.
refine (bind ξvars = manyFresh _  $\Sigma \xi_0$ ; _).
apply FreshMon.
destruct ξvars as [vars ξpf].
destruct ξpf as [ξ pf].
refine (cnd ξ vars _ >>>= fun it => _).
apply FreshMon.
auto.
refine (return OK _).
exists ξ.

```

```
apply (ileaf it).  
Defined.
```

```
End HaskProofToStrong.
```